

# Joint Routing and Charging Problem of Electric Vehicles

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# Outline

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- 1. Research Background**
- 2. Electric Vehicles Routing Problem**
- 3. Electric Vehicles Routing Problem with Time Flexibility**
- 4. Conclusion**

# Outline

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## 1. Research Background

## 2. Electric Vehicles Routing Problem

## 3. Electric Vehicles Routing Problem with Time Flexibility

## 4. Conclusion

# Research background

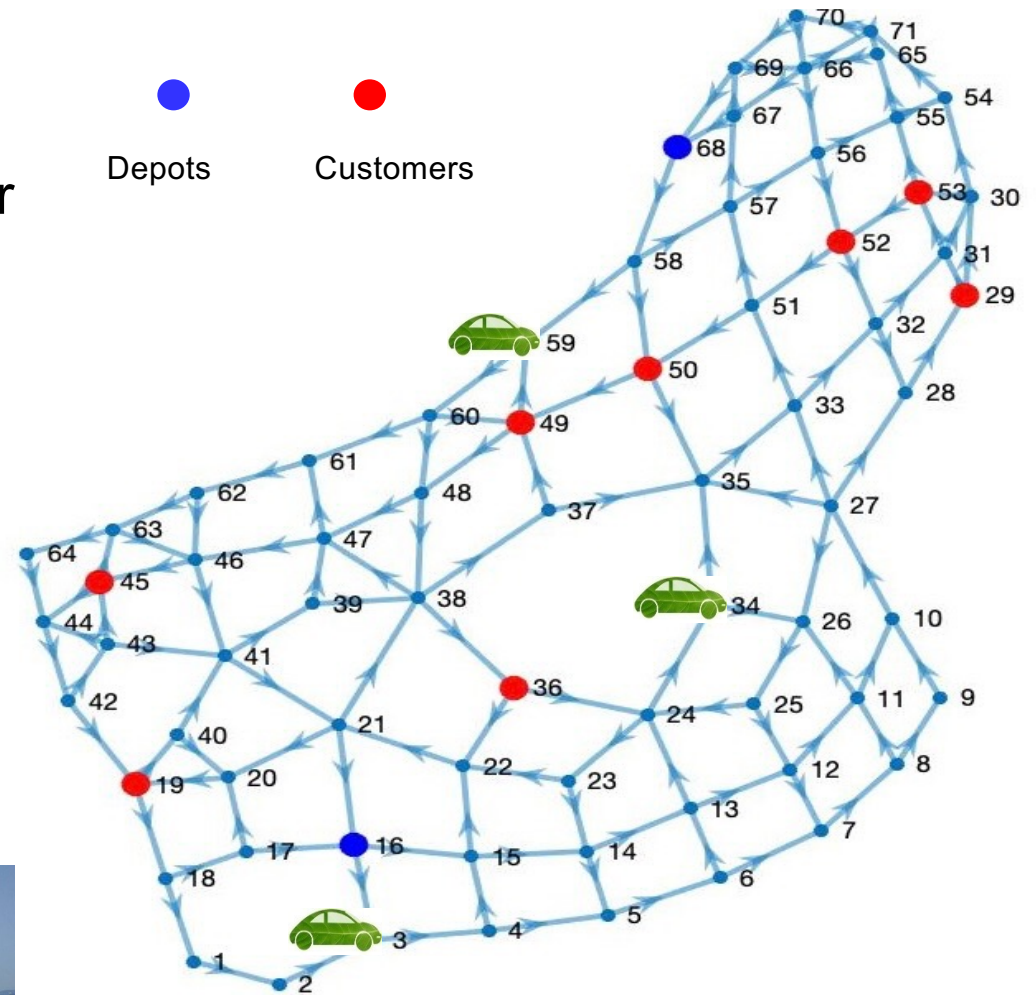
## □ Transportation Problem

### ➤ Definition

- Design the optimal routes for a fleet of vehicles, to serve the customers.

### ➤ Applications

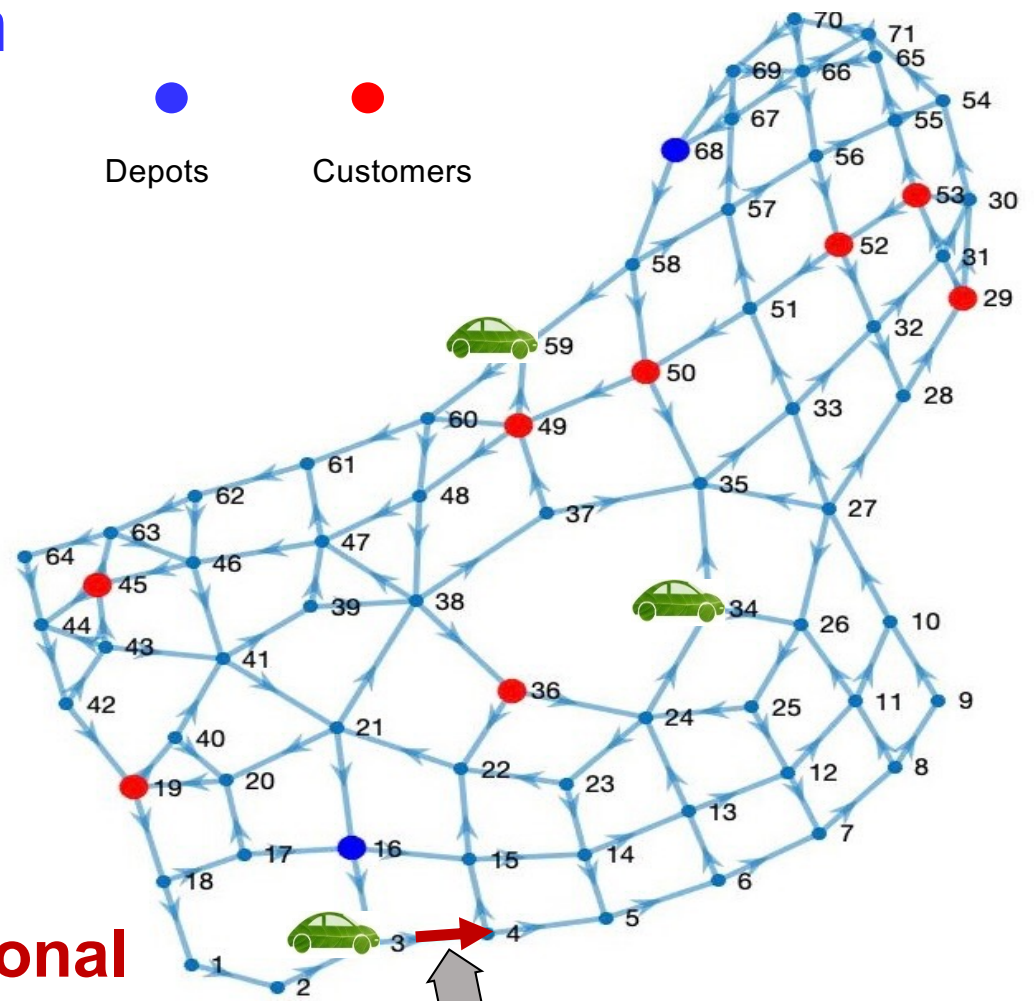
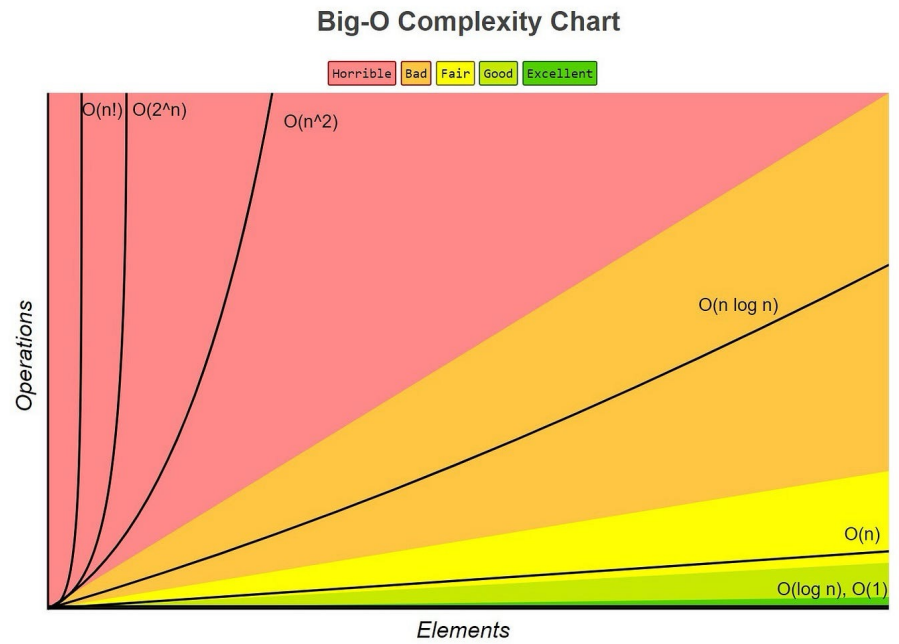
- Package delivery, and taxi operation



**Vehicle routing problem**

# Research background

## Vehicle routing problem



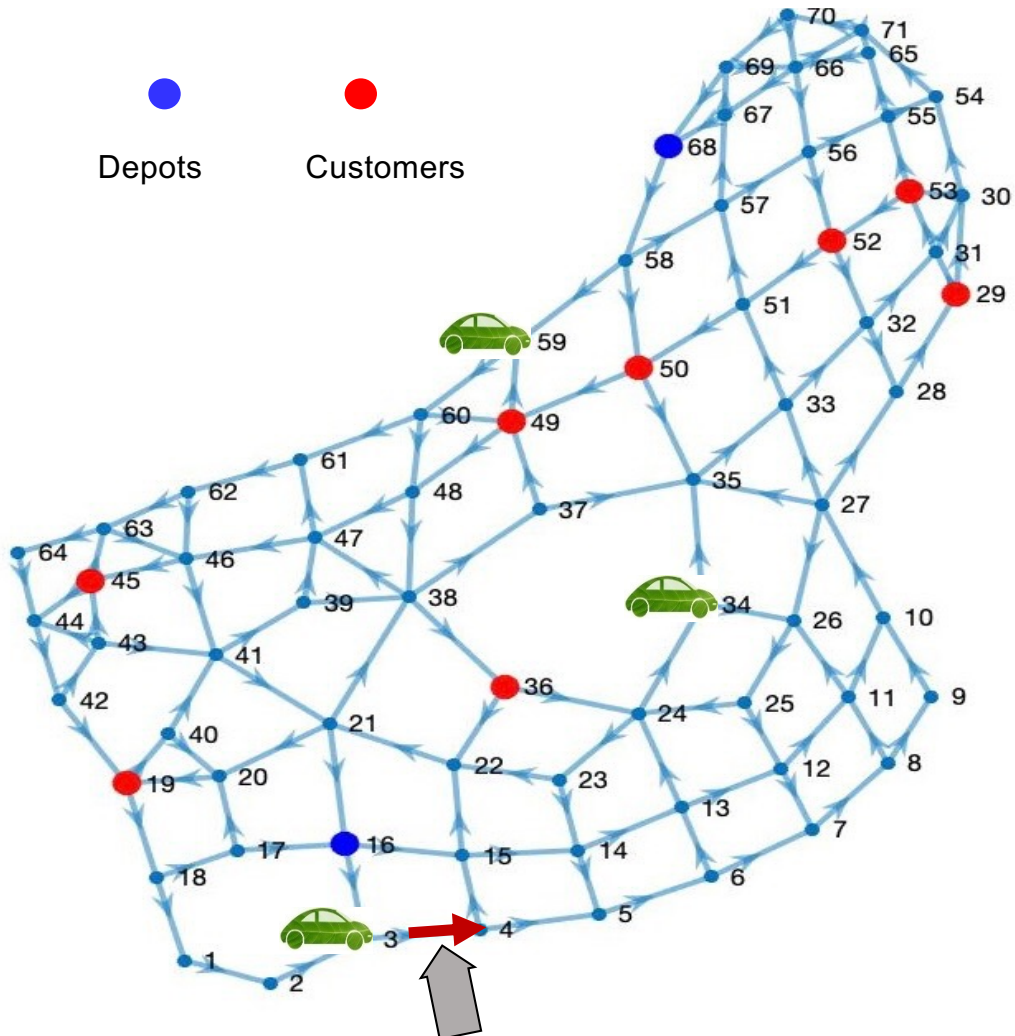
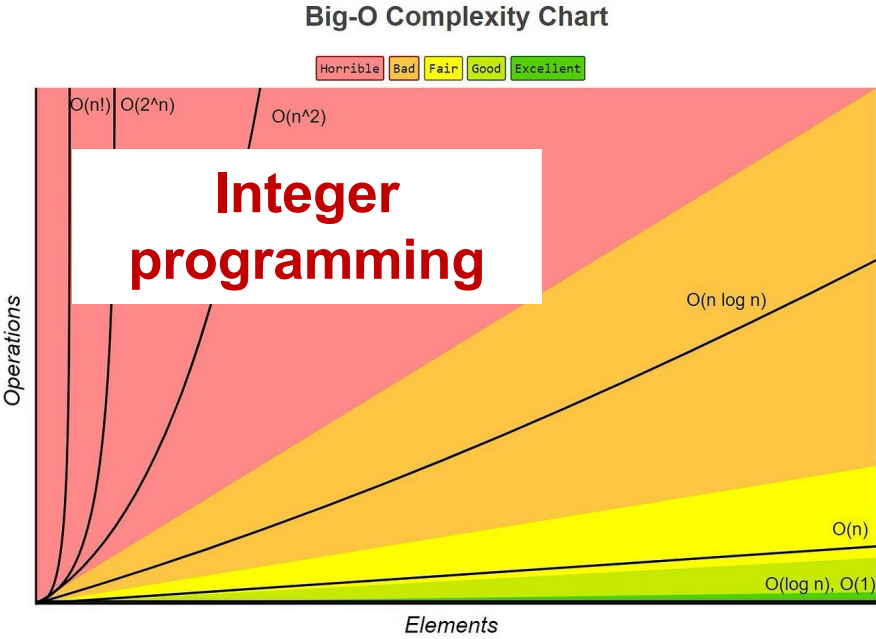
**High computational complexity!**

$$x_{ij}^k \in \{0,1\}$$

[1] Toth, Paolo, and Daniele Vigo. Vehicle routing: problems, methods, and applications. SIAM, 2014.

# Research background

## Vehicle routing problem



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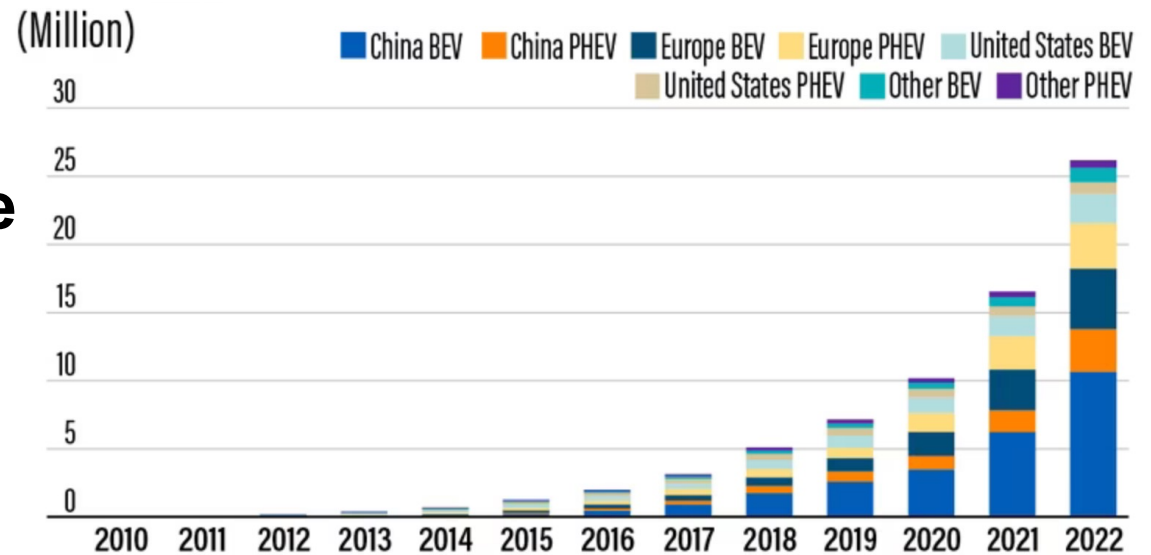
# Research background

## □ Electric vehicle

- Lower emissions
- Reduced dependence on fossil resources

## □ Limitations

- Range anxiety
- Long charging time



Vehicle routing  
problem



Electric vehicle routing  
problem







# Research background

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## □ Literature review

### ➤ Exact algorithms

- Branch and bound method (TITS'18)
- Branch, cut and price method (IJPR'22, TS22 )
- Solver-based method (TSG'18, TVT'20 )

**High computation complexity**

### ➤ Approximate approaches

- Dual decomposition (TSG'18, TVT'20 )
- ADMM (TSG'18, TVT'20 )

**Suboptimal, high computation complexity**

### ➤ Heuristics methods

- Variable neighborhood search (TSG'18, TVT'20 )
- Genetic algorithm (TSG'18, TVT'20 )
- Deep reinforcement learning method (TSG'18, TVT'20 )

**Suboptimal**

# Research background

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## □ Motivation

### ➤ Large-sized electric vehicle routing problem

- The computation time of large-sized EVRP, is quite long!
- The online EVRP requires a computation-efficient algorithm!

### ➤ Collaborative scheduling of electric vehicle fleet and customers

- Fixed pickup time of customers reduces the efficiency of electric transportation system.
- To characterize the collaborative scheduling problem, how to formulate a suitable mathematical model?

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4. Conclusion

# EV routing problem

## □ Mathematical model (MIP)

### ➤ Objective function

$$\min_{x_{ij}^k, r_i^k} \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{V}} \sum_{j \in \mathcal{V}} \left( c_i + \omega_T T_{ij} + r_i^k p_i + \omega_T r_i^k g_i \right) x_{ij}^k$$

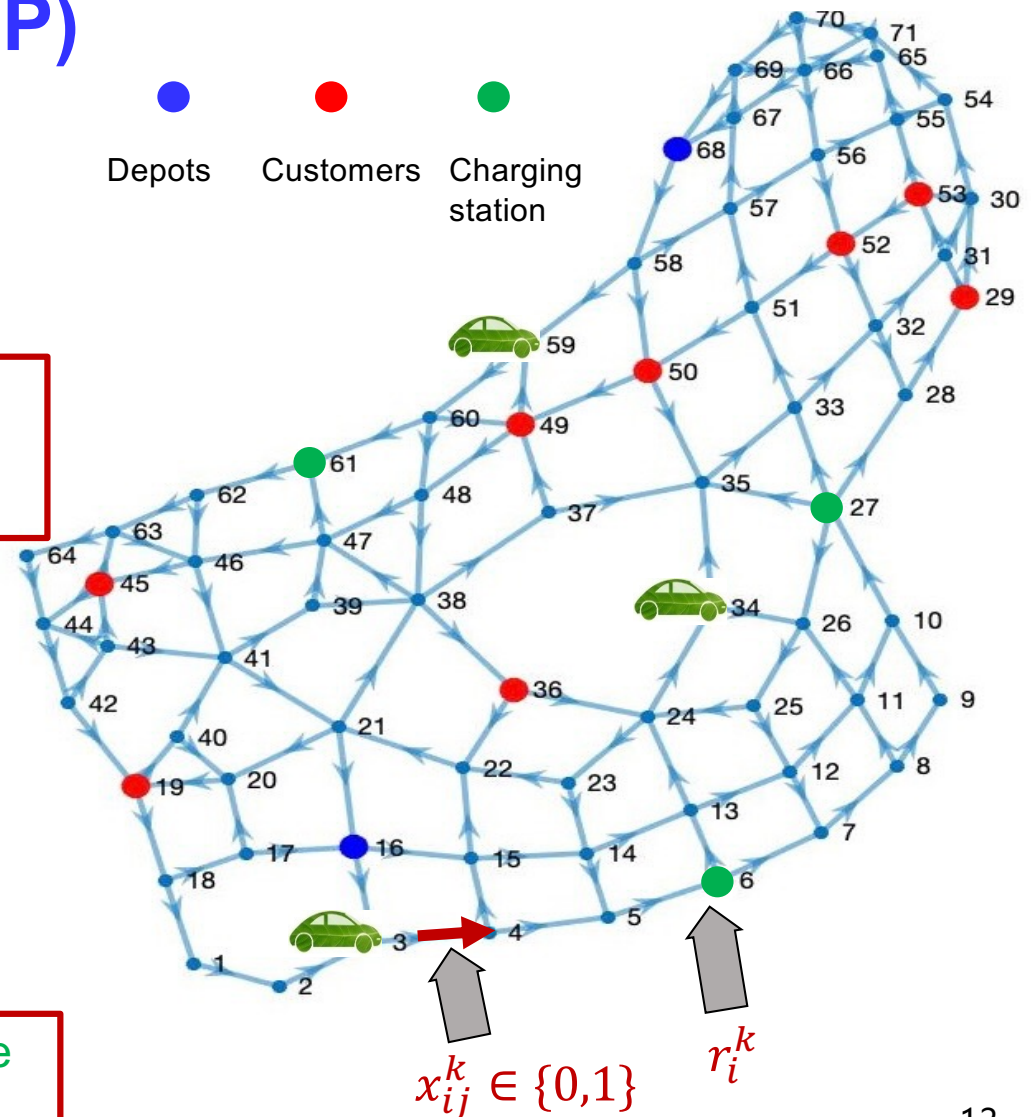
Operation cost = usage cost + travel time  
+ charging cost  
+ charging time

### ➤ Constraints

- Vehicle flow

$$\sum_{j \in \mathcal{V}} x_{ij}^k - \sum_{j \in \mathcal{V}} x_{ji}^k = b_i, b_{v_1} = 1, b_{v_n} = -1, b_i = 0,$$

Starting from the start depot, after visiting the customers, return to the end depot



# EV routing problem

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## □ Mathematical model (MIP)

### ➤ Constraints

- Visiting constraint

$$\sum_{k \in \mathcal{K}} \sum_{j \in \mathcal{V}} x_{ij}^k \leq 1,$$

Customer cannot be visited by more than one EV

- Visiting time

$$t_j \geq (T_{ij} + g_i r_i^k + t_i) x_{ij}^k,$$

The visiting time of customers

- SoC of EV

$$E_j^k = \sum_{i \in \mathcal{V}} (E_i^k + r_i^k - e_{ij}) x_{ij}^k,$$
$$0 \leq E_j^k \leq \bar{E}.$$

The state of charge of EV

# EV routing problem

## □ Mathematical model (MIP)

$$\min_{x_{ij}^k \in \mathbb{B}, r_i^k \in \mathbb{R}} \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{V}} \sum_{j \in \mathcal{V}} (c_i + \omega_T T_{ij} + r_i^k p_i + \omega_T r_i^k g_i) x_{ij}^k$$

$$\text{s.t.} \quad \sum_{j \in \mathcal{V}} x_{ij}^k - \sum_{j \in \mathcal{V}} x_{ji}^k = b_i, \quad \forall i \in \mathcal{V}, k \in \mathcal{K}, b_{v_1} = 1, b_{v_n} = -1, b_i = 0$$

$$\sum_{k \in \mathcal{K}} \sum_{j \in \mathcal{V}} x_{ij}^k \leq 1, \quad \forall i \in \mathcal{R}$$

$$t_j \geq (T_{ij} + g_i r_i^k + t_i) x_{ij}^k, \quad \forall i \in \mathcal{V} \setminus v_n, j \in \mathcal{V} \setminus v_1, k \in \mathcal{K}$$

$$E_j^k = \sum_{i \in \mathcal{V}} (E_i^k + r_i^k - e_{ij}) x_{ij}^k, \quad \forall j \in \mathcal{V} \setminus v_1, k \in \mathcal{K},$$

$$0 \leq E_j^k \leq \bar{E}, \quad \forall j \in \mathcal{V} \setminus v_1, k \in \mathcal{K}.$$





# EV routing problem

## □ Mathematical model (MIP)

$$\min_{x_{ij}^k \in \mathbb{B}, r_i^k \in \mathbb{R}} \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{V}} \sum_{j \in \mathcal{V}} (c_i + \omega_T T_{ij} \quad \text{[redacted]}) x_{ij}^k$$

$$\text{s.t.} \quad \sum_{j \in \mathcal{V}} x_{ij}^k - \sum_{j \in \mathcal{V}} x_{ji}^k = b_i, \quad \forall i \in \mathcal{V}, k \in \mathcal{K}, b_{v_1} = 1, b_{v_n} = -1, b_i = 0$$

$$\sum_{k \in \mathcal{K}} \sum_{j \in \mathcal{V}} x_{ij}^k \leq 1, \quad \forall i \in \mathcal{R}$$

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Theorem 1 (**Dimension reduction**)



# EV routing problem

## □ Two stage method (Routing problem)

### ➤ Elimination of bilinear terms

- Eliminating bilinear term  $r_i x_{ij}$  ( $r_i$  is replaced with  $e_{ij}$ )

### ➤ Exact LP relaxation

- Vehicle flow constraints. are totally unimodular, satisfying sufficient condition of exact LP relaxation
- Time constraint does not break LP relaxation exactness

IP  LP

$$\min_{x_{ij} \in \{0,1\}} \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{V}} \sum_{j \in \mathcal{V}} (c_i + \omega_T T_{ij}) x_{ij}$$

$$\text{s.t. } t_j \geq (T_{ij} + t_i + g_i r_i) x_{ij}, \forall i \in \mathcal{V} \setminus v_n, j \in \mathcal{V} \setminus v_1$$

$$\sum_{j \in \mathcal{V}} x_{ij} - \sum_{j \in \mathcal{V}} x_{ij} = b_i, \quad \forall i \in \mathcal{V},$$

$$b_{v_1} = 1, b_{v_n} = -1, b_i = 0$$

$$\sum_{k \in \mathcal{K}} \sum_{j \in \mathcal{V}} x_{ij} \leq 1, \quad \forall i \in \mathcal{R}$$



# EV routing problem

## □ Two stage method (Charging scheduling)

➤ Given  $x_{ij}$ , obtain the path  $P_k$  for each EV

How to get a better solution?

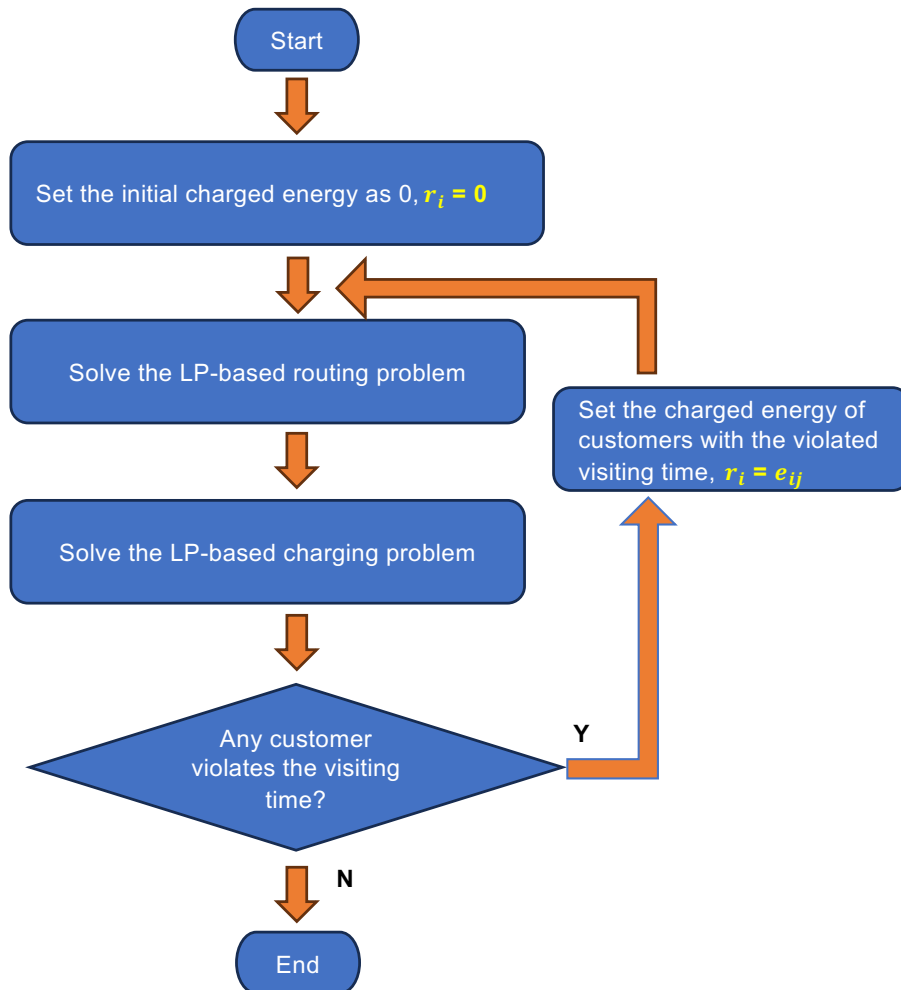
$$\begin{aligned} \min_{r_i \in R, i \in P_k} \quad & \omega_T \sum_{i \in P_k} r_i g_i + \sum_{i \in P_k} r_i p_i \\ \text{s.t.} \quad & E_j = E_i + r_i - e_{ij}, \forall j \in P_k \setminus v_1 \\ & 0 \leq E_j^k \leq \bar{E} \end{aligned}$$

The charging scheduling problem is a **LP** problem



# EV routing problem

## Iterative two stage method



## Performance guarantee

### Two stage method

- Both routing and charging problem are **LP problems**, which reduces the computation time

### Iterative method

- Converge in finite iterations

### Optimality gap

$$C_{MIP} \leq C_{ILP} \leq C_{TLP}$$

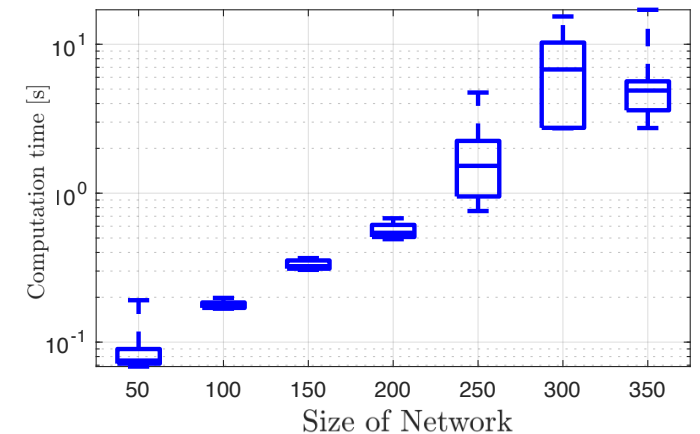
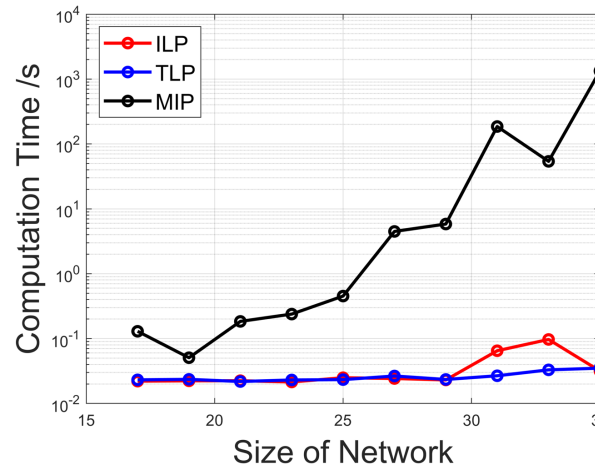
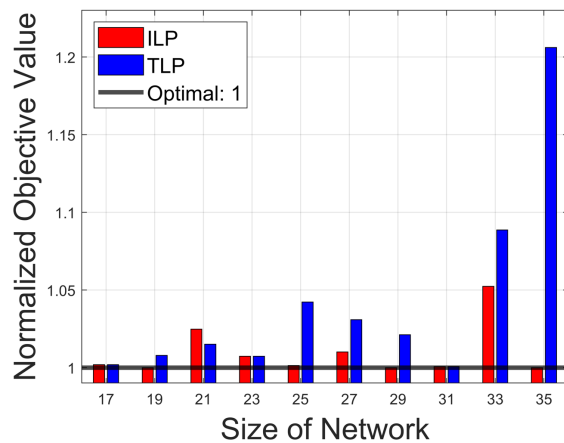
### Computation time

$$T_{MIP} \gg T_{ILP} \geq T_{TLP}$$

# EV routing problem

## Simulation results

- Comparison between TLP, and ILP
- Large-sized problem



The optimality gap of ILP lower than 10%  
The optimality gap of TLP lower than 20%



The computation time of ILP 0.1 s  
The computation time of TLP 0.05 s

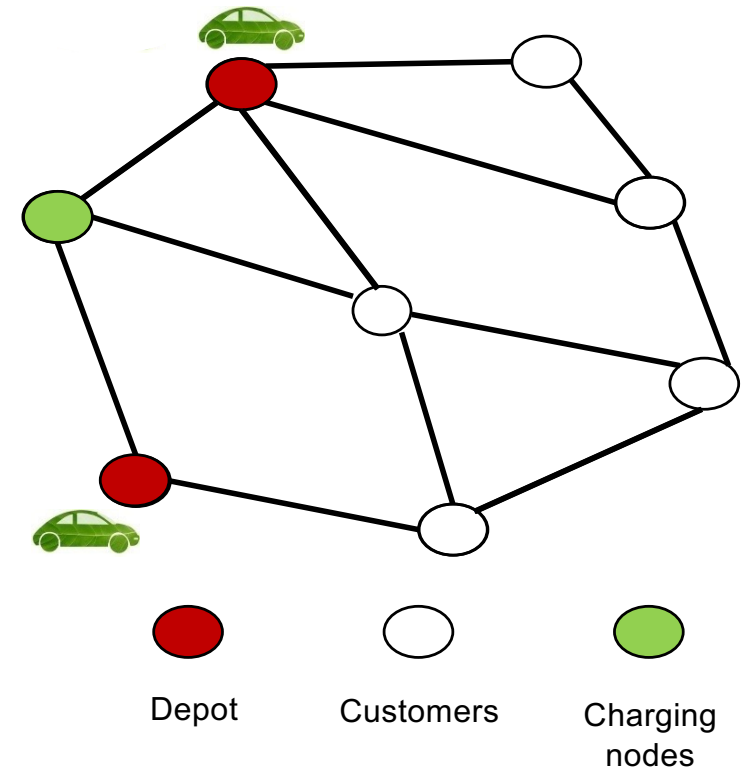
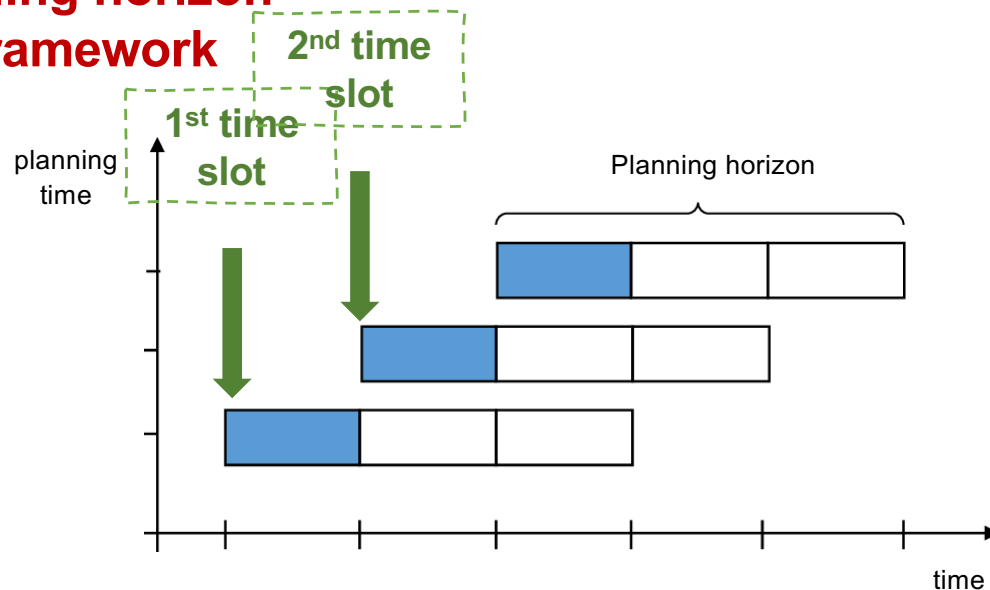
350 nodes and  
35 EVs solved  
in 10 seconds

# Online EV routing problem

## □ Online routing problem

- Real-time generated customers
- The uncertain charging prices

### Rolling horizon framework



Online EVRP

Rolling horizon

Multidepot routing

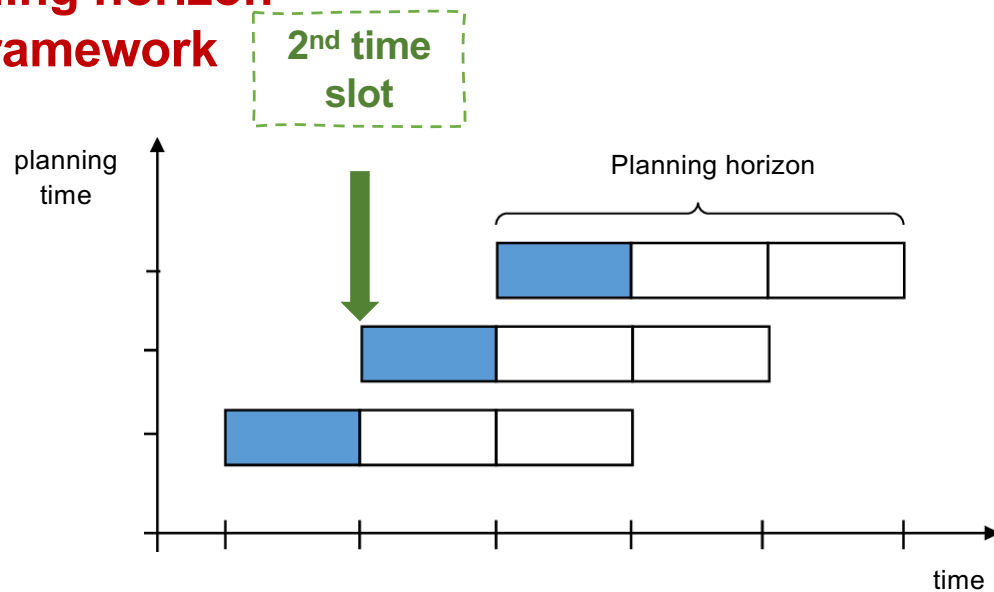


# Online EV routing problem

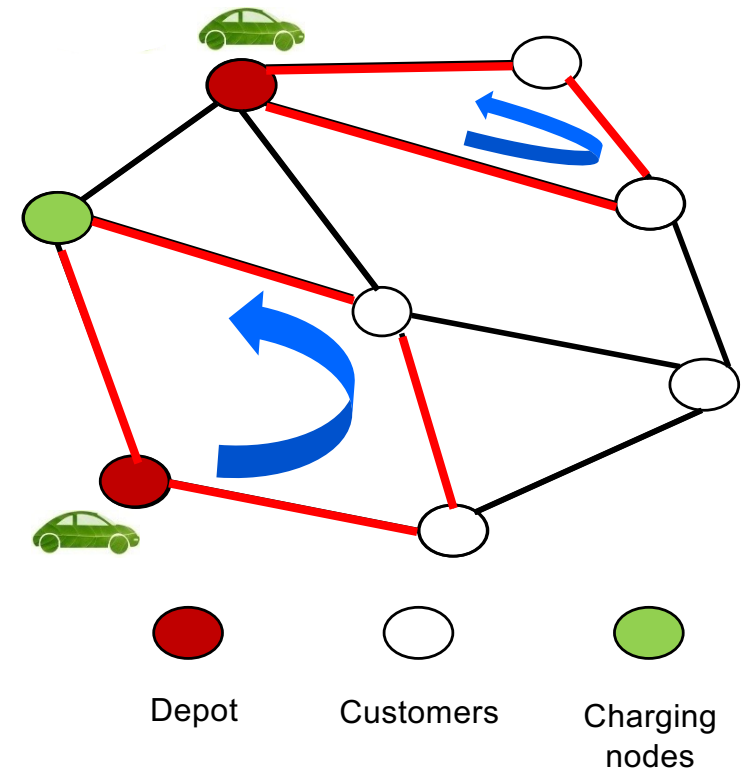
## □ Online routing problem

- Real-time generated customers
- The uncertain charging prices

### Rolling horizon framework



Challenges: Multi-depot routing problem



Online EVRP

Rolling horizon

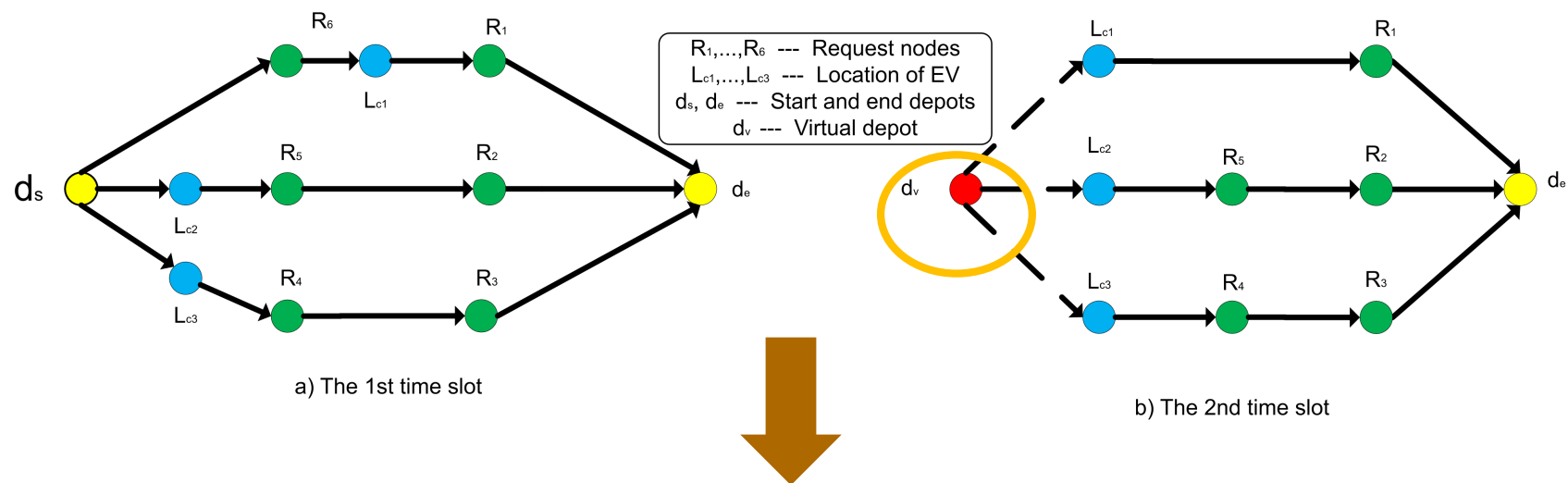
Multidepot routing

# Online EV routing problem

## Multi-depot routing problem

### ➤ The equivalent transformation rule

- The *road selection variable and travel time* of virtual roads are set as **1,0**, respectively.
- The energy consumption of virtual roads is set as **the difference between the battery capacity and current SoC** of EV.

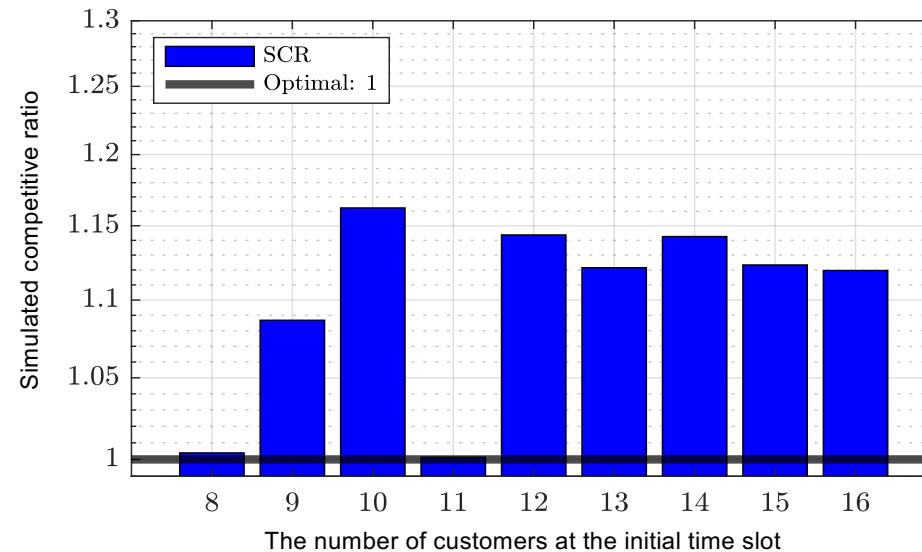


**Homogenous EV at the virtual depot**

# Online EV routing problem

## □ Simulation results

### ➤ Evaluation of the simulated competitive ratio



The SCR of ILP is lower than **1.2**, demonstrating that the ILP works well in online routing problem

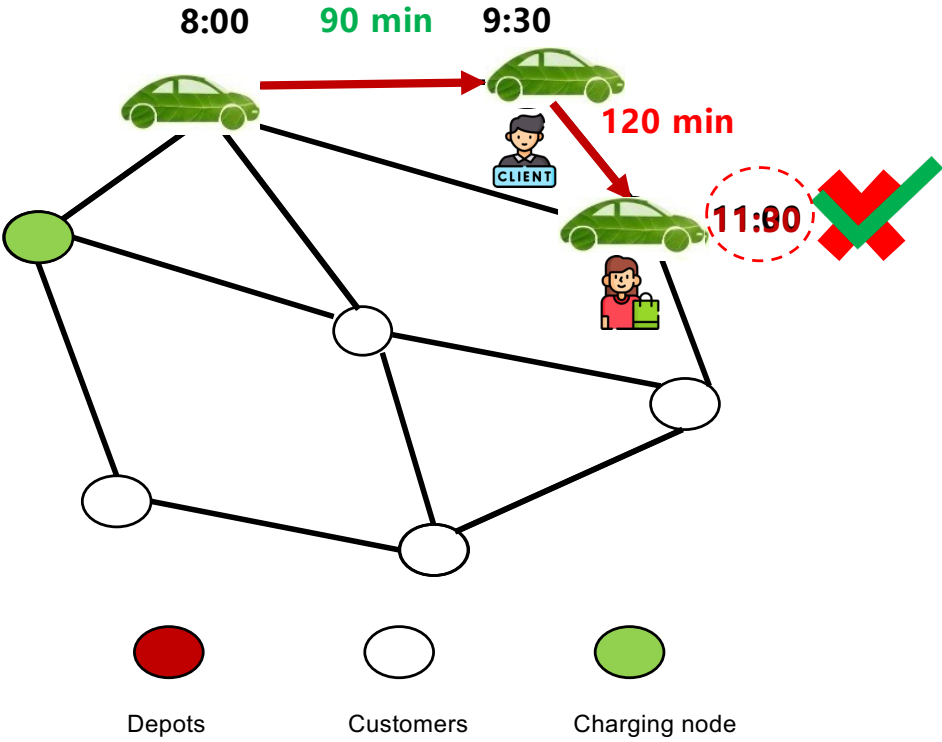
# Outline

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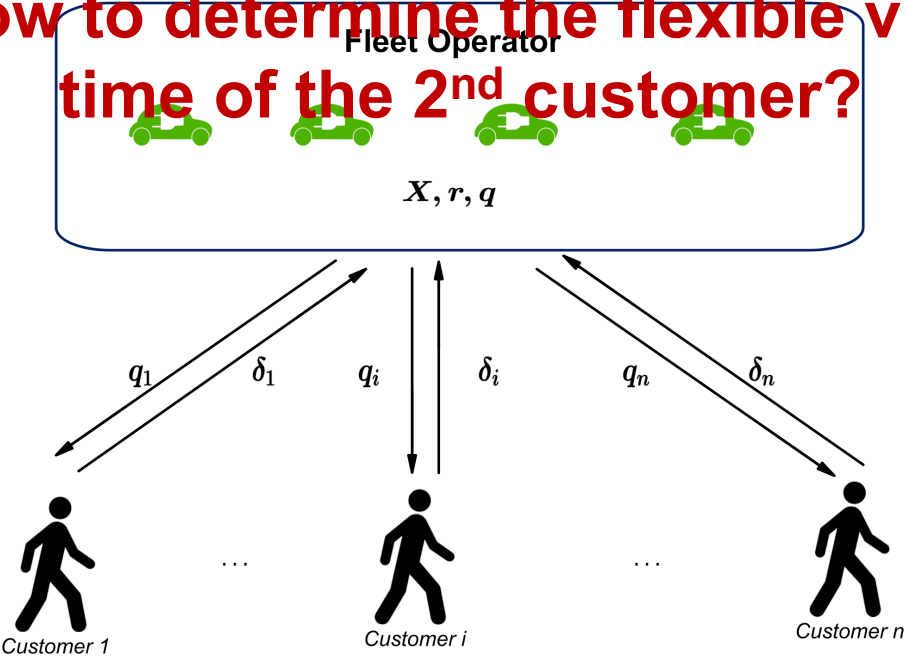
1. Research Background
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4. Conclusion

# EVRP with time flexibility

## □ Time flexibility



How to determine the flexible visit time of the 2<sup>nd</sup> customer?



# EVRP with time flexibility

## □ Mathematical model of the fleet operator

### ➤ Objective function

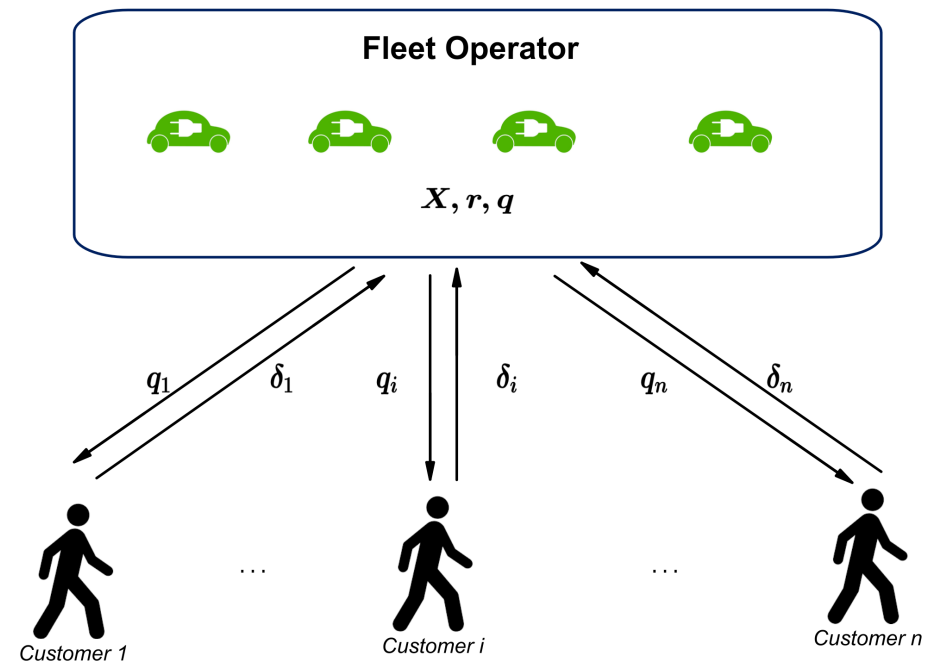
$$\min_{x_{ij}^k \in \mathbb{B}, r_i^k, q_j, t_j \in \mathbb{R}} \sum_{i \in \mathcal{C}} \sum_{\tau \in \Lambda} \frac{p_{i\tau} B_{i\tau} \Delta \tau}{g_i} + \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{V}} \sum_{j \in \mathcal{V}} (c_i + \omega_T T_{ij} + \omega_T r_i^k g_i) x_{ij}^k + \sum_{j \in \mathcal{R}} q_j \delta_j^* \sum_{i \in \mathcal{V}} \sum_{k \in \mathcal{K}} x_{ij}^k$$

Operation cost of the fleet +  
delivery fee discount

### ➤ Constraint

- SoC
- Vehicle flow
- Visit time
- Flexible time window

$$t_j^L \leq t_j \leq t_j^L + \delta_j, \quad \forall j \in \mathcal{R}.$$





# EVRP with time flexibility

## Mathematical model of customers

$$\min_{\delta_j \in \mathbb{R}} \mathcal{J}(\delta_j) - q_j \delta_j,$$

$$\text{s.t. } 0 \leq \delta_j \leq \bar{\delta}_j$$

Minimize the inconvenience cost +  
maximize the discount of delivery fee

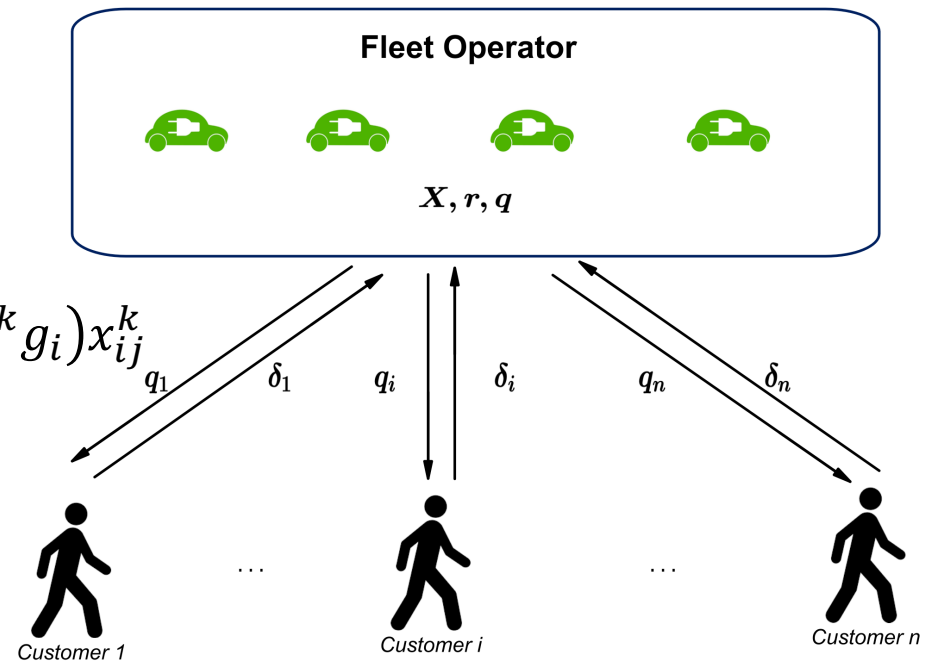
## The joint model of the fleet and customers

$$\min_X \sum_{i \in \mathcal{C}} \sum_{\tau \in \Lambda} \frac{p_{i\tau} B_{i\tau} \Delta \tau}{g_i} + \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{V}} \sum_{j \in \mathcal{V}} (c_i + \omega_T T_{ij} + \omega_T r_i^k g_i) x_{ij}^k$$

$$+ \sum_{j \in \mathcal{R}} q_j \delta_j^* \sum_{i \in \mathcal{V}} \sum_{k \in \mathcal{K}} x_{ij}^k$$

$$\text{s.t. } \delta_j^* \in \arg \min_{\delta_j} \{ \mathcal{J}(\delta_j) - q_j \delta_j, 0 \leq \delta_j \leq \bar{\delta}_j \},$$

$$\forall j \in \mathcal{R},$$



How to solve the joint optimization problem?

# EVRP with time flexibility

## □ The equivalent reformulation method

### ➤ KKT optimality condition

$$\nabla \mathcal{J}(\delta_j) - q_j - \sigma_j + u_j = 0, \forall j \in \mathcal{R}$$

$$0 \leq \bar{\delta}_j - \delta_j \leq M\psi_j^1, \forall j \in \mathcal{R}$$

$$0 \leq u_j \leq M(1 - \psi_j^1), \forall j \in \mathcal{R}$$

$$0 \leq \delta_j \leq M\psi_j^2, \forall j \in \mathcal{R}$$

$$0 \leq \sigma_j \leq M(1 - \psi_j^2), \forall j \in \mathcal{R}$$

Stationary  
condition

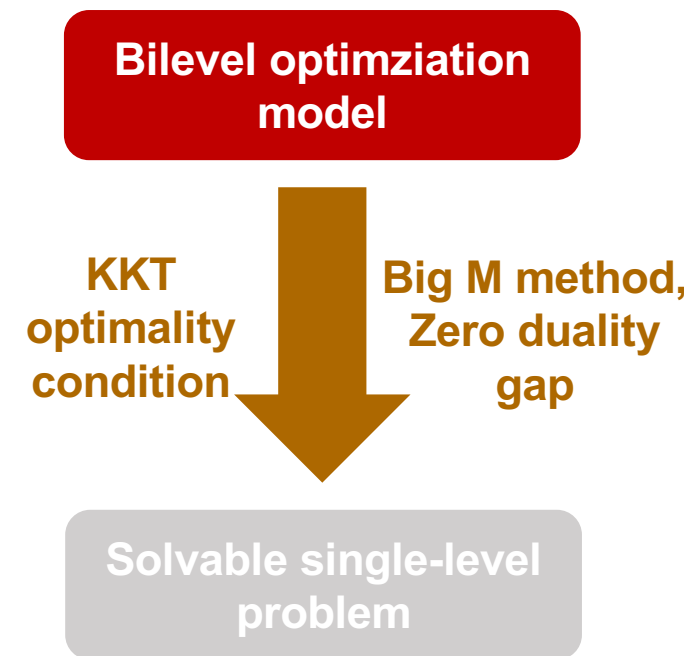
Linearized  
complementarity  
condition

### ➤ Linearization of nonlinear terms

$$\eta_{ijk}^1 \geq \omega_T r_i^k g_i - M(1 - x_{ij}^k), \forall i \in \mathcal{V}, j \in \mathcal{V}, k \in \mathcal{K}$$

$$\eta_j^2 \geq \mathcal{J}(\delta_j) + u_j \bar{\delta}_j - \phi^*(\delta_j^*) - M \left( 1 - \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{J}} x_{ij}^k \right), \forall j \in \mathcal{R}.$$

Zero duality gap of the  
lower-level problem

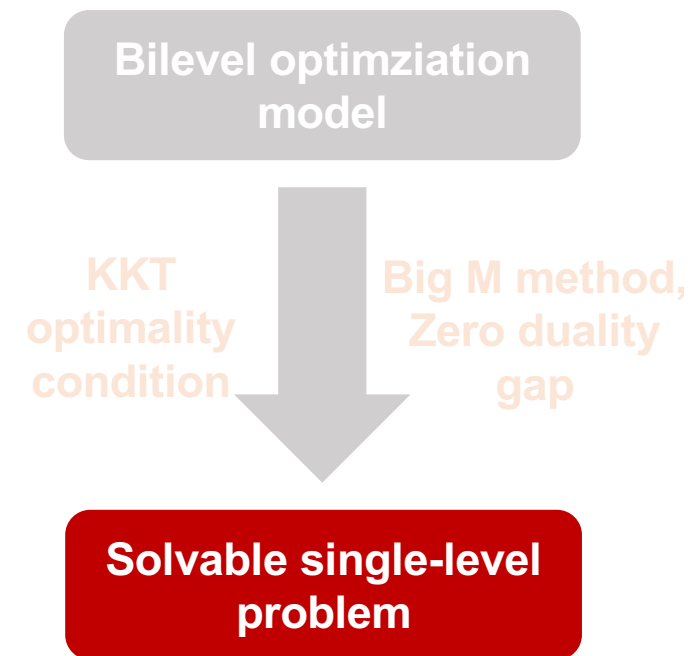
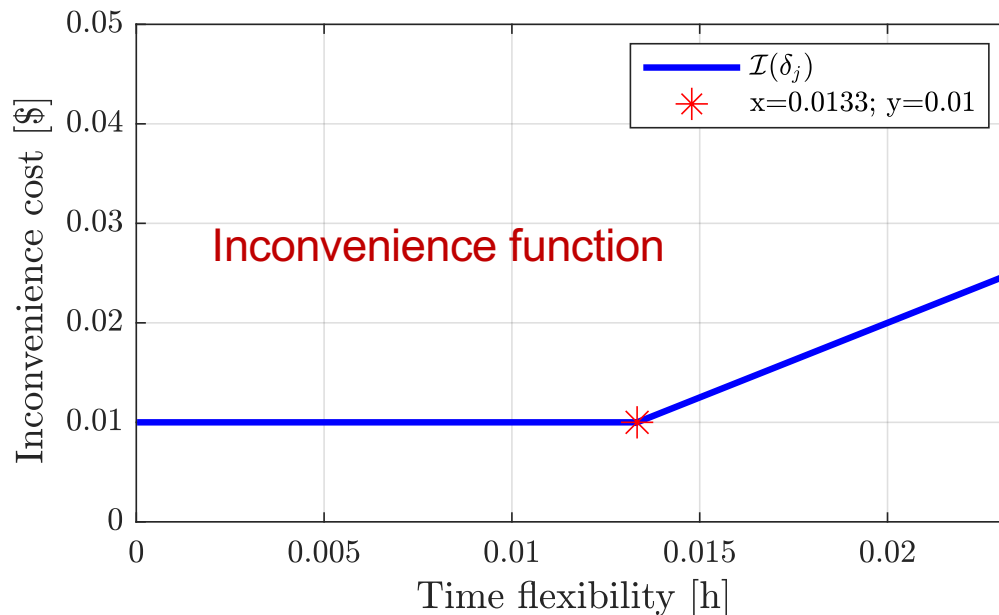


# EVRP with time flexibility

## □ Linearized single-level problem

- MIP-based single-level Problem
- Solved by commercial solver

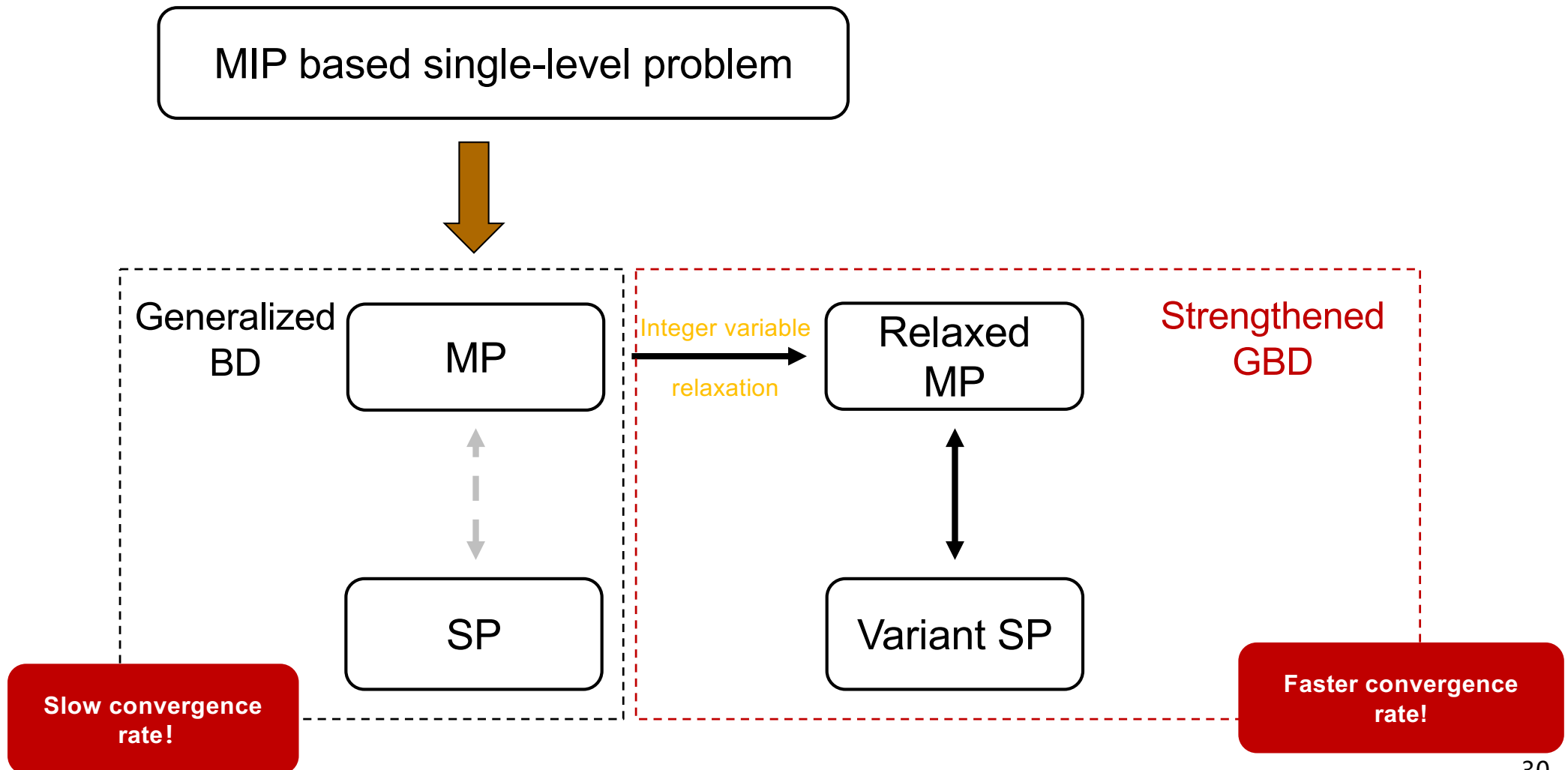
$$\min_x \sum_{k \in \mathcal{K}} \sum_{i \in \mathcal{V}} \sum_{j \in \mathcal{V}} (\eta_{ijk}^1 + \eta_j^2 + (\omega_T T_{ij} + c_i) x_{ij}^k)$$



**How to design the distributed algorithm to protect the privacy of both sides?**

# EVRP with time flexibility

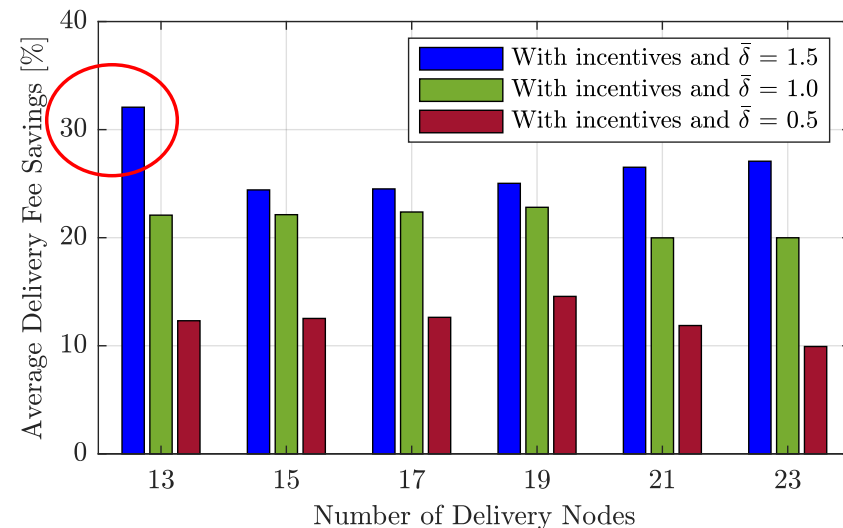
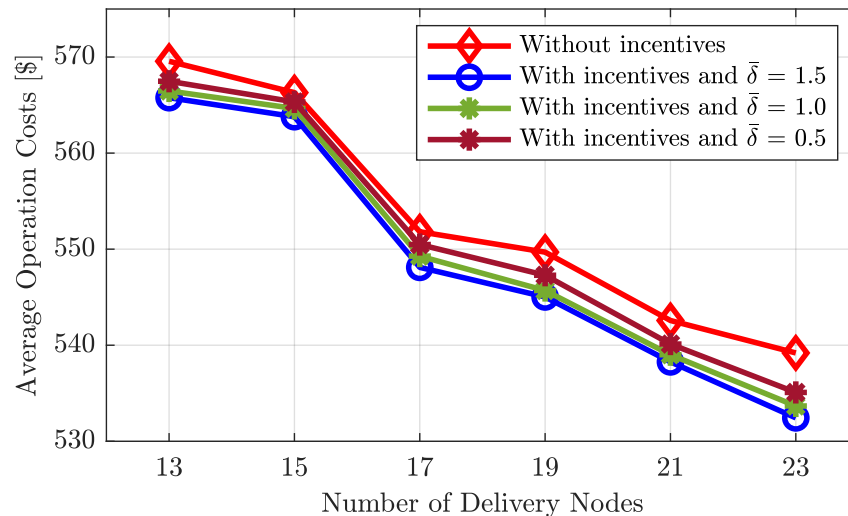
## □ The algorithm details



# EVRP with time flexibility

## Simulation results

### ➤ Impact of Different Time-flexibility Levels $\bar{\delta}_j$

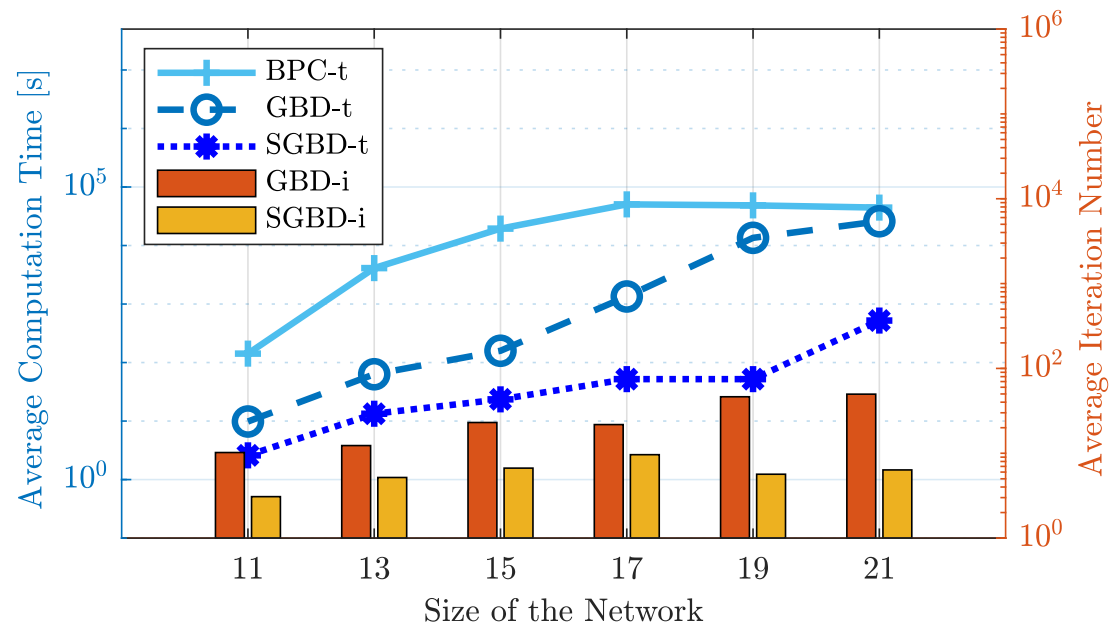


The operation cost becomes **smaller** with increasing values of  $\bar{\delta}_j$ ;  
The delivery fee saving becomes **larger** with increasing values of  $\bar{\delta}_j$

# EVRP with time flexibility

## Simulation results

### Comparison results between SGBD, GBD, BCP



SGBD shows a better performance than GBD, and BCP in terms of **the number of iterations and computation time**



# EVRP with time flexibility

## □ Simulation results

### ➤ The scalability of SGBD

Instance	$ \mathcal{V}  = 11$ $ \mathcal{E}  = 22$	$ \mathcal{V}  = 51$ $ \mathcal{E}  = 102$	$ \mathcal{V}  = 101$ $ \mathcal{E}  = 202$	$ \mathcal{V}  = 151$ $ \mathcal{E}  = 302$
Run time (s)	2.06	121.47	837.35	*
$\epsilon$	0	0.41%	2.33%	12.23%

SGBD can solve the EVRP with **up to 150 nodes and 15 EVs**

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# Conclusion

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- ❑ Design the approximate method to solve large-sized EVRP **in polynomial time**
- ❑ Devise the **rolling horizon based** method to solve the online EVRP with the near-optimal solution
- ❑ Propose **a bilevel optimization model**, to characterize the mathematical model of customers and the electric vehicle fleet
- ❑ Design the KKT condition based equivalent reformulation method, and the **strengthened generalized Benders decomposition method**



***Thank you!***



***Q & A***

## **Acknowledgement**

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